

# **Austinian truth in Martin-Löf type theory**

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# Austinian truth

Barwise and Perry (*Situations and Attitudes*) paraphrase Austin:

a statement is true when the actual situation to which it refers is of the type described by the statement.

Austin's original ('Truth', 1961)

A statement is said to be true when the historic state of affairs to which it is correlated by the demonstrative conventions (the ones to which it 'refers') is of a type with which the sentence used in making it is correlated by the descriptive conventions.

- Statements are true of something (part of the world) not true *simpliciter*
- Truth *simpliciter* can be derived: there is some part of the world of which the statement is true
- Statements are correlated with types

# Martin-Löf Type Theory

- Objects are of types
- Propositions are regarded as types of proofs (“propositions as types” principle)
- Proofs are objects, e.g. the proofs of *there is a prime number between 212 and 222* are the prime numbers between 212 and 222

# Proofs of non-mathematical propositions

Ranta in *Type-Theoretical Grammar* draws a parallel with Davidson's event-based approach:

A proof of

*Amundsen flew over the North Pole*

is

a flight by Amundsen over the North Pole

In terms of situation theory:

A proof of

*Amundsen flew over the North Pole*

is

a **situation** in which Amundsen flies over the North Pole

Infons (states of affairs, soas) – types of situations

<b>TT</b>	<b>ST</b>
types of proofs (propositions)	infons

### **Sentences of situation theory**

$s \models \sigma$  (“situation  $s$  supports infon (soa)  $\sigma$ ” or “ $s$  is of type  $\sigma$ ”)

### **Situation theoretic objects**

$(s \models \sigma)$ , an Austinian proposition, true just in case  $s \models \sigma$

# Judgements

## Type theoretical judgements

$a : T$  (“object  $a$  is of type  $T$ ”)

## Type theoretical objects

an object  $a$ , of type  $T - a : T$

No objects corresponding to judgements as such

# Judgements concerning proofs

$p : \text{see}(a, b)$

$p$  is a proof that  $a$  sees  $b$

context:  $a:\text{Ind}, b:\text{Ind}$

# **A note on individuals**

What is the type Ind?

Schemes of individuation

# Records

Contexts are not objects but records are.

We shall use record types which have fields corresponding both to the context and the propositional judgement.

$$T = \left[ \begin{array}{l} x \quad : \text{Ind} \\ y \quad : \text{Ind} \\ \text{prf} : \text{see}(x,y) \end{array} \right]$$

$$\left[ \begin{array}{l} x \quad = a \\ y \quad = b \\ \text{prf} = p \\ \dots \end{array} \right] : T \text{ just in case } a : \text{Ind}, b : \text{Ind} \text{ and } p : \text{see}(a, b)$$

# Why records?

Three reasons independent of situation theory:

- allow us to reify contexts and quantify over them
- correspond to discourse representation structures – a record type is inhabited (“true”) just in case there is something of the appropriate type corresponding to each of its labels.
- records can be values in fields as in typed feature structures – useful for compositional semantics, parallels with HPSG

## Situations as records

Records and record types give us the possibility of grouping together collections of infon-like objects.

$$\left[ \begin{array}{l} s : \left[ \begin{array}{l} x : \text{Ind} \\ y : \text{Ind} \\ s_1 : \text{see}(x,y) \\ s_2 : \text{see}(y,x) \end{array} \right] \end{array} \right]$$

$$\left[ \begin{array}{l} x : \text{Ind} \\ y : \text{Ind} \\ s : \left[ \begin{array}{l} s_1 : \text{see}(x,y) \\ s_2 : \text{see}(y,x) \end{array} \right] \end{array} \right]$$

## Constraints on universes of records

*smoke means fire*

Any universe of records which obeys this constraint is such that if for some labels  $l$  and  $s$  it contains a record of type

$$\left[ \begin{array}{l} l : \text{Loc} \\ s : \text{smoke}(l) \end{array} \right]$$

then it also, for some labels  $l'$  and  $s'$ , contains a record of type

$$\left[ \begin{array}{l} l' : \text{Loc} \\ s' : \text{fire}(l') \end{array} \right]$$

*kissing involves touching*

Applies to single situations (you cannot perceive *a* kiss *b* without perceiving *a* touch *b*)

Any record in the universe which for some labels  $x,y$  and  $s$  is of type

$$\left[ \begin{array}{l} x : \text{Ind} \\ y : \text{Ind} \\ s : \text{kiss}(x,y) \end{array} \right]$$

is also, for some label  $s'$ , of type

$$\left[ \begin{array}{l} x : \text{Ind} \\ y : \text{Ind} \\ s' : \text{touch}(x,y) \end{array} \right]$$

Note that  $x$  and  $y$  in the second record to depend on the first record.

## Can constraints be internalized?

$$\left[ f : \left( \left[ \begin{array}{l} l : \text{Loc} \\ s : \text{smoke}(l) \end{array} \right] \right) \left[ \begin{array}{l} l' : \text{Loc} \\ s' : \text{fire}(l') \end{array} \right] \right]$$

$$\left[ f : \left( r : \left[ \begin{array}{l} x : \text{Ind} \\ y : \text{Ind} \\ s : \text{kiss}(x,y) \end{array} \right] \right) \left[ \begin{array}{l} r' : \left[ \begin{array}{l} s' : \text{touch}(r.x,r.y) \end{array} \right] \\ c : \text{eq}(\text{RecType}, r', r) \end{array} \right] \right]$$

Functions correspond to universal quantification.

What to do if the constraint is soft or *fallible*? Functions in TT are not normally partial.

## Channel types

$$\mathcal{C} = \left[ \begin{array}{l} \text{signal} : T_1 \\ \text{target} : T_2 \end{array} \right]$$

A record  $r : \mathcal{C}$  is a *link classified by  $\mathcal{C}$* .

An agent attuned to  $\mathcal{C}$  on observing an object of type  $T_1$  will predict (defeasibly) the existence of an object of type  $T_2$ .

According to  $\mathcal{C}$  an object  $a : T_1$  *carries the information that  $T_2$  is inhabited*.

## Dependent target types

Becomes of some interest when we consider that the target type may depend on the signal type.

$$\left[ \begin{array}{l} \text{signal} : \left[ \begin{array}{l} x : \text{Ind} \\ y : \text{Ind} \\ c_1 : r(x,y) \end{array} \right] \\ \text{target} : \left[ \begin{array}{l} c_2 : r'(\text{signal}.x, \text{signal}.y) \end{array} \right] \end{array} \right]$$

Note that we do not know what the target type is until we have determined a particular signal.

# Records and attitudes

## *A simple treatment of the attitudes*

Record types as the object of the attitude (corresponding to a ST treatment where attitudes are relations between individuals and in-fons).

$$\left[ \begin{array}{l} x : \text{Ind} \\ c_1 : \text{girl}(x) \\ \\ c_2 : \text{believe}(x, \left[ \begin{array}{l} y : \text{Ind} \\ c_3 : \text{man}(y) \\ z : \text{Ind} \\ c_4 : \text{donkey}(z) \\ c_5 : \text{own}(y, z) \end{array} \right] ) \end{array} \right]$$

- The object of belief is Austinian.
- Austinian propositions could be represented as pairs of records and record types (coded as a record).

# Perception attitudes

*a man sees that a donkey kicked a farmer*

$$\left[ \begin{array}{l} x : \text{Ind} \\ c_1 : \text{man}(x) \\ c_5 : \text{see}(x, \left[ \begin{array}{l} y : \text{Ind} \\ c_2 : \text{donkey}(y) \\ z : \text{Ind} \\ c_3 : \text{farmer}(z) \\ c_4 : \text{kick}(y, z) \end{array} \right] ) \end{array} \right]$$

*Veridicality*

$$\left[ f : (r: \left[ \begin{array}{l} x : \text{Ind} \\ T : \text{RecType} \\ c : \text{see}(x, T) \end{array} \right]) \left[ a : r.T \right] \right]$$

Note that this does not require that the man saw a donkey-kicking-farmer situation – he may just have seen donkey hoof marks on the farmer’s legs.

# Naked infinitive perception complements

*a man sees a donkey kick a farmer*

$$\left[ \begin{array}{l} x : \text{Ind} \\ c_1 : \text{man}(x) \\ s : \left[ \begin{array}{l} y : \text{Ind} \\ c_2 : \text{donkey}(y) \\ z : \text{Ind} \\ c_3 : \text{farmer}(z) \\ c_4 : \text{kick}(y, z) \end{array} \right] \\ c_5 : \text{see}(x,s) \end{array} \right]$$

Veridicality is required.

Also perception of the donkey-kicking-farmer situation.

# Types are not enough

- Proper names
- Presuppositions

*Sandy kicks a farmer*

$$\left[ \begin{array}{l} y : \text{Ind} \\ c_2 : \text{named}(y, \text{“Sandy”}) \\ z : \text{Ind} \\ c_3 : \text{farmer}(z) \\ c_4 : \text{kick}(y, z) \end{array} \right]$$

## Frames of mind

A *family of types* – function from records of a given type to a record type.

$$\lambda r : \left[ \begin{array}{l} y : \text{Ind} \\ c_2 : \text{named}(y, \text{“Sandy”}) \end{array} \right] \left[ \begin{array}{l} z : \text{Ind} \\ c_3 : \text{farmer}(z) \\ c_4 : \text{kick}(r.y, z) \end{array} \right]$$

A *setting* for a frame of mind is a record in its domain.

The *presupposition* associated with a frame of mind is the type characterizing its domain.

A *mental state* is a pair consisting of a frame of mind (agent internal) and a setting (agent external).

The *content of a mental state* is the result of applying the frame of mind to the setting.

Advantage over Cooper and Ginzburg (1996) : no restricted objects

# Pierre

## *Frame of mind*

$$\lambda r : \left[ \begin{array}{l} x : \text{Ind} \\ c_1 : \text{named}(x, \text{"Londres"}) \\ y : \text{Ind} \\ c_2 : \text{named}(y, \text{"London"}) \end{array} \right] \left[ \begin{array}{l} c_3 : \text{pretty}(r.x) \\ c_4 : \neg\text{pretty}(r.y) \end{array} \right]$$

## *Setting*

$$\left[ \begin{array}{l} x = \text{london} \\ c_1 = \text{pf named}(\text{london} \text{"Londres"}) \\ y = \text{london} \\ c_2 = \text{pf named}(\text{london} \text{"London"}) \end{array} \right]$$

## *Content of Pierre's mental state*

$$\left[ \begin{array}{l} c_3 : \text{pretty}(\text{london}) \\ c_4 : \neg\text{pretty}(\text{london}) \end{array} \right]$$

# Compositionality

Sample derivation: *every man owns a donkey*

*a donkey*

$$\lambda R_1:([x:\text{Ind}])\text{RecType } \lambda R_2:([x:\text{Ind}])\text{RecType} \left[ \begin{array}{l} \text{par} \quad : \quad [x : \text{Ind}] \\ \text{restr} \quad : \quad R_1 @ \text{par} \\ \text{scope} \quad : \quad R_2 @ \text{par} \end{array} \right]$$

@

$$\lambda r:[x:\text{Ind}]([c:\text{donkey}(r.x)])$$

=

$$\lambda R_2:([x:\text{Ind}])\text{RecType} \left[ \begin{array}{l} \text{par} \quad : \quad [x : \text{Ind}] \\ \text{restr} \quad : \quad [c : \text{donkey}(\text{par}.x)] \\ \text{scope} \quad : \quad R_2 @ \text{par} \end{array} \right]$$

*own a donkey*

$\lambda \mathcal{N} : (([x:\text{Ind}])\text{RecType})\text{RecType} \lambda r_1 : [x:\text{Ind}] (\mathcal{N} @ \lambda r_2 : [x:\text{Ind}] ([c:\text{own}(r_1.x, r_2.x)] @$

$\lambda R_2 : ([x:\text{Ind}])\text{RecType} \left[ \begin{array}{l} \text{par} \quad : \quad [x : \text{Ind}] \\ \text{restr} \quad : \quad [c : \text{donkey}(\text{par}.x)] \\ \text{scope} \quad : \quad R_2 @ \text{par} \end{array} \right]$

=

$\lambda r_1 : [x:\text{Ind}] \left( \left[ \begin{array}{l} \text{par} \quad : \quad [x : \text{Ind}] \\ \text{restr} \quad : \quad [c : \text{donkey}(\text{par}.x)] \\ \text{scope} \quad : \quad [c : \text{own}(r_1.x, \text{par}.x)] \end{array} \right] \right)$

*every man*

$\lambda R_1:([\mathbf{x}:\text{Ind}])\text{RecType } \lambda R_2:([\mathbf{x}:\text{Ind}])\text{RecType}$

$\left[ \mathbf{f} : (r : \left[ \begin{array}{l} \text{par} : [\mathbf{x} : \text{Ind}] \\ \text{restr} : R_1 @ \text{par} \end{array} \right]) \right] R$

@

$\lambda r:[\mathbf{x}:\text{Ind}](\left[ \mathbf{c}:\text{man}(r.\mathbf{x}) \right])$

=

$\lambda R_2:([\mathbf{x}:\text{Ind}])\text{RecType}$

$\left[ \mathbf{f} : (r : \left[ \begin{array}{l} \text{par} : [\mathbf{x} : \text{Ind}] \\ \text{restr} : \left[ \mathbf{c} : \text{man}(\text{par}.\mathbf{x}) \right] \end{array} \right]) \right] R_2 @ r.\text{par}$

*every man owns a donkey*

$\lambda R_2:([x:\text{Ind}])\text{RecType}$

$$\left[ \text{f} : (r : \left[ \begin{array}{l} \text{par} : [x : \text{Ind}] \\ \text{restr} : [c : \text{man}(\text{par}.x)] \end{array} \right]) R_2 @ r.\text{par} \right]$$

@

$$\lambda r_1:[x:\text{Ind}] \left( \left[ \begin{array}{l} \text{par} : [x : \text{Ind}] \\ \text{restr} : [c : \text{donkey}(\text{par}.x)] \\ \text{scope} : [c : \text{own}(r_1.x,\text{par}.x)] \end{array} \right] \right)$$

=

$$\left[ \text{f} : (r : \left[ \begin{array}{l} \text{par} : [x : \text{Ind}] \\ \text{restr} : [c : \text{man}(\text{par}.x)] \end{array} \right]) \left[ \begin{array}{l} \text{par} : [x : \text{Ind}] \\ \text{restr} : [c : \text{donkey}(\text{par}.x)] \\ \text{scope} : [c : \text{own}(r.\text{par}.x,p)] \end{array} \right] \right]$$

### *Flattening*

$$\left[ f : (r : \left[ \begin{array}{l} \text{par.x} : \text{Ind} \\ \text{restr.c} : \text{man}(\text{par.x}) \end{array} \right]) \left[ \begin{array}{l} \text{par.x} : \text{Ind} \\ \text{restr.c} : \text{donkey}(\text{par.x}) \\ \text{scope.c} : \text{own}(r.\text{par.x},\text{par.x}) \end{array} \right] \right]$$

### *Relabelling*

$$\left[ f : (r : \left[ \begin{array}{l} x : \text{Ind} \\ c_1 : \text{man}(x) \end{array} \right]) \left[ \begin{array}{l} y : \text{Ind} \\ c_2 : \text{donkey}(y) \\ c_3 : \text{own}(r.x,y) \end{array} \right] \right]$$

### *r-suppression (syntactic sugar)*

$$\left[ f : \left( \left[ \begin{array}{l} x : \text{Ind} \\ c_1 : \text{man}(x) \end{array} \right] \right) \left[ \begin{array}{l} y : \text{Ind} \\ c_2 : \text{donkey}(y) \\ c_3 : \text{own}(x,y) \end{array} \right] \right]$$

$$\llbracket [\mathbf{V}_S \text{ believes}] \rrbracket = \lambda T:\text{RecType } \lambda r:[\mathbf{x}:\text{Ind}]([\mathbf{c}:\text{believe}(r.\mathbf{x},T)])$$

$$\llbracket [\mathbf{V}_S \text{ sees}] \rrbracket = \lambda T:\text{RecType } \lambda r:[\mathbf{x}:\text{Ind}]([\mathbf{c}:\text{see}_t(r.\mathbf{x},T)])$$

$$\begin{aligned} \llbracket [\mathbf{V}_{NI} \text{ sees}] \rrbracket = & \lambda \mathcal{N}:(([\mathbf{x}:\text{Ind}])\text{RecType})\text{RecType} \\ & \lambda R:([\mathbf{x}:\text{Ind}])\text{RecType} \\ & \lambda r:[\mathbf{x}:\text{Ind}] \left( \begin{array}{l} \mathbf{s} : \mathcal{N} @ R \\ \mathbf{c} : \text{see}_s(r.\mathbf{x},\mathbf{s}) \end{array} \right) \end{aligned}$$

$$\llbracket [\mathbf{VP} \mathbf{V}_S \mathbf{S}] \rrbracket = \llbracket [\mathbf{V}_S] \rrbracket @ \llbracket [\mathbf{S}] \rrbracket$$

$$\llbracket [\mathbf{VP} \mathbf{V}_{NI} \mathbf{NP} \mathbf{VP}[\text{inf}]] \rrbracket = \llbracket [\mathbf{V}_{NI}] \rrbracket @ \llbracket [\mathbf{NP}] \rrbracket @ \llbracket [\mathbf{VP}[\text{inf}]] \rrbracket$$

# Records as resources

Domain restrictions

*every/the man owns a donkey*

Resource situations in situation semantics

“everything which is a man in situation  $s$ ”

“everything which is a man in record  $r$ ”

We use

$$r \models T$$

to represent the type of objects in record  $r$  which are of type  $T$

$a : r \models T$  just in case for some  $\ell, r : [\ell:T]$  and  $r.\ell = a$

*Anchored resources (specific)*

$\lambda r_1 : \text{Rec}$

$\lambda r_2 : \text{Rec}$

$$\left[ \begin{array}{l} \mathbf{f} : (r : \left[ \begin{array}{l} \mathbf{x} : r_1 \models \text{Ind} \\ \mathbf{c}_1 : r_1 \models \text{man}(\mathbf{x}) \end{array} \right]) \left[ \begin{array}{l} \mathbf{y} : r_2 \models \text{Ind} \\ \mathbf{c}_2 : r_2 \models \text{donkey}(\mathbf{y}) \\ \mathbf{c}_3 : \text{own}(r.\mathbf{x},\mathbf{y}) \end{array} \right] \end{array} \right] \\ @ \text{res}_1 @ \text{res}_2$$

“Every man in  $\text{res}_1$  owns a donkey in  $\text{res}_2$ .”

*Non-anchored resources (non-specific, existentially quantified)*

$$\left[ \begin{array}{l} \text{res}_1 : \text{Rec} \\ \text{res}_2 : \text{Rec} \\ \mathbf{f} : (r : \left[ \begin{array}{l} \mathbf{x} : \text{res}_1 \models \text{Ind} \\ \mathbf{c}_1 : \text{res}_1 \models \text{man}(\mathbf{x}) \end{array} \right]) \left[ \begin{array}{l} \mathbf{y} : \text{res}_2 \models \text{Ind} \\ \mathbf{c}_2 : \text{res}_2 \models \text{donkey}(\mathbf{y}) \\ \mathbf{c}_3 : \text{own}(r.\mathbf{x},\mathbf{y}) \end{array} \right] \end{array} \right]$$

“There are resources  $\text{res}_1, \text{res}_2$  such that every man in  $\text{res}_1$  owns a donkey in  $\text{res}_2$ .”

*Quantification over resources (a kind of generic)*

$$\left[ f : (r : \left[ \begin{array}{l} \text{res}_1 : \text{Rec} \\ x : \text{res}_1 \models \text{Ind} \\ c_1 : \text{res}_1 \models \text{man}(x) \end{array} \right]) \left[ \begin{array}{l} \text{res}_2 : \text{Rec} \\ y : \text{res}_2 \models \text{Ind} \\ c_2 : \text{res}_2 \models \text{donkey}(y) \\ c_3 : \text{own}(r.x,y) \end{array} \right] \right]$$

“For every resource  $\text{res}_1$  and every man  $x$  in  $\text{res}_1$  there is a resource  $\text{res}_2$  such that  $x$  owns a donkey in  $\text{res}_2$ .”

## Records (situations) as modules

Types restricted to resources import information from the resource (cf importing from modules in programming languages).

If  $r : [ a : r_1 \models T ]$   
then  $r : [ a : T ]$

$r$  can be used as a resource for another record.

This kind of “information spreading” not present in situation theory.

# Conclusions

Records in Martin-Löf type theory

- allow us to combine various natural language “technologies”
  - Montague lambda technology
  - DRT, dynamic semantics
  - typed feature structures (HPSG)
- give us a promising approach to situations and intensionality, including
  - an Austinian approach to perception and attitude complements
  - constraints/channels
- exploit key features of MLTT
  - dependent types
  - propositions as types