

# Using Type Theory with Records for HPSG

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# The Records and Dialogue Semantics Project

<http://www.ling.gu.se/~cooper/records/> (These slides are there)

- Records, types and computational dialogue semantics

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CHALMERS

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Computing Science

- Computer Science and Engineering – Chalmers University of Technology and Göteborg University

- Robin Cooper, Thierry Coquand, Staffan Larsson, Peter Ljunglöf, Bengt Nordström, Arne Ranta

# Four linguistic theories + dialogue management

- Montague semantics
- DRT
- situation semantics
- HPSG
- issue based dialogue management (Larsson)

Main advantage: you can get aspects of all five theories going at the same time.

# HPSG in type theory with records

Collaboration with Jonathan Ginzburg, King's College, London.

cf. Ginzburg's paper at HPSG in Leuven

# Contents

<b>1</b>	<b>Records in type theory</b>	<b>6</b>
<b>2</b>	<b>Records as linguistic objects</b>	<b>11</b>
<b>3</b>	<b>A case where we don't seem to need unification</b>	<b>20</b>
<b>4</b>	<b>Putting content into TTR-HPSG</b>	<b>29</b>
<b>5</b>	<b>A case where we seem to need not to have unification</b>	<b>34</b>
<b>6</b>	<b>Conclusions</b>	<b>43</b>
<b>A</b>	<b>Abstract and concrete syntax</b>	<b>44</b>
⇐	contents	5/53

# 1. Records in type theory

## Records and record types

If  $a_1 : T_1, a_2 : T_2(a_1), \dots, a_n : T_n(a_1, a_2, \dots, a_{n-1})$ ,  
the record:

$$\left[ \begin{array}{l} l_1 = a_1 \\ l_2 = a_2 \\ \dots \\ l_n = a_n \\ \dots \end{array} \right]$$

is of type:

$$\left[ \begin{array}{l} l_1 : T_1 \\ l_2 : T_2(l_1) \\ \dots \\ l_n : T_n(l_1, l_2, \dots, l_{n-1}) \end{array} \right]$$

*a man owns a donkey*

Record type:

$$\left[ \begin{array}{l} x : Ind \\ c_1 : \text{man}(x) \\ y : Ind \\ c_2 : \text{donkey}(y) \\ c_3 : \text{own}(x,y) \end{array} \right]$$

Record:

$$\left[ \begin{array}{l} x = a \\ c_1 = p_1 \\ y = b \\ c_2 = p_2 \\ c_3 = p_3 \end{array} \right]$$

where

$a, b$  are of type *Ind*, individuals

$p_1$  is a proof of  $\text{man}(a)$

$p_2$  is a proof of  $\text{donkey}(b)$

$p_3$  is a proof of  $\text{own}(a, b)$

⇐ contents



*a man owns a donkey*

*Content (intension)* is a record type:

$$\left[ \begin{array}{l} x : Ind \\ c_1 : \text{man}(x) \\ y : Ind \\ c_2 : \text{donkey}(y) \\ c_3 : \text{own}(x,y) \end{array} \right]$$

- a record of this type may have additional fields
- the types  $\text{man}(x)$ ,  $\text{donkey}(y)$ ,  $\text{own}(x,y)$  are dependent types of proofs

# Manifest fields

Coquand, Pollack and Takeyama

If  $a : T$ , then  $T_a$  is a *singleton type*

$b : T_a$  iff  $b = a$

A manifest field in a record type is one whose type is a singleton type, e.g.

$[ x : T_a ]$

written for convenience as

$[ x=a : T ]$

Allows record types to be “progressively instantiated”.

We will allow dependent unique types, i.e. where  $a$  can be represented by a path in a record type.

## 2. Records as linguistic objects

*man*

$$\left[ \begin{array}{l} \text{phon} = [\text{man}] \\ \text{cat} = \text{n} \\ \text{agr} = \left[ \begin{array}{l} \text{num} = \text{sg} \\ \text{gen} = \text{masc} \\ \text{pers} = \text{third} \end{array} \right] \end{array} \right]$$

*fish*

$$\left[ \begin{array}{l} \text{phon} = [\text{fish}] \\ \text{cat} = \text{n} \\ \text{agr} = \left[ \begin{array}{l} \text{num} = \text{sg} \\ \text{gen} = \text{neut} \\ \text{pers} = \text{third} \end{array} \right] \end{array} \right]$$

$$\left[ \begin{array}{l} \text{phon} = [\text{fish}] \\ \text{cat} = \text{n} \\ \text{agr} = \left[ \begin{array}{l} \text{num} = \text{pl} \\ \text{gen} = \text{neut} \\ \text{pers} = \text{third} \end{array} \right] \end{array} \right]$$

$$\left[ \begin{array}{l} \text{phon} = [\text{fish}] \\ \text{cat} = \text{n} \\ \text{agr} = \left[ \begin{array}{l} \text{num} = \text{sg} \\ \text{gen} = \text{masc} \\ \text{pers} = \text{third} \end{array} \right] \end{array} \right]$$

$$\left[ \begin{array}{l} \text{phon} = [\text{fish}] \\ \text{cat} = \text{n} \\ \text{agr} = \left[ \begin{array}{l} \text{num} = \text{sg} \\ \text{gen} = \text{fem} \\ \text{pers} = \text{third} \end{array} \right] \end{array} \right]$$

...

## Part of a type system (ignoring semantics)

$$Noun \equiv \left[ \begin{array}{l} \text{phon} : Phon \\ \text{cat=n} : Cat \\ \text{agr} : \left[ \begin{array}{l} \text{num} : Number \\ \text{gen} : Gender \\ \text{pers} : Person \end{array} \right] \end{array} \right]$$

*Phon*  $\equiv$  [*Lex*]

the, a, fish, man, ... : *Lex*

n, det, np, v, vp, s, ... : *Cat*

sg, pl : *Number*

masc, fem, neut : *Gender*

first, second, third : *Person*

[⇐ contents](#)

$$Det \equiv \left[ \begin{array}{l} \text{phon} : \textit{Phon} \\ \text{cat=det} : \textit{Cat} \\ \text{agr} : \left[ \begin{array}{l} \text{num} : \textit{Number} \\ \text{gen} : \textit{Gender} \\ \text{pers} : \textit{Person} \end{array} \right] \end{array} \right]$$

$$Agr \equiv \left[ \begin{array}{l} \text{num} : \textit{Number} \\ \text{gen} : \textit{Gender} \\ \text{pers} : \textit{Person} \end{array} \right]$$



## A singleton (fully specified) type

<table style="border-collapse: collapse; width: 100%;"> <tr> <td style="padding: 5px;">phon=[man]</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Phon</i></td> </tr> <tr> <td style="padding: 5px;">cat=n</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Cat</i></td> </tr> <tr> <td style="padding: 5px;">agr</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"> <table style="border-collapse: collapse; width: 100%;"> <tr> <td style="padding: 5px;">num=sg</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Number</i></td> </tr> <tr> <td style="padding: 5px;">gen=masc</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Gender</i></td> </tr> <tr> <td style="padding: 5px;">pers=third</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Person</i></td> </tr> </table> </td> </tr> </table>	phon=[man]	:	<i>Phon</i>	cat=n	:	<i>Cat</i>	agr	:	<table style="border-collapse: collapse; width: 100%;"> <tr> <td style="padding: 5px;">num=sg</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Number</i></td> </tr> <tr> <td style="padding: 5px;">gen=masc</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Gender</i></td> </tr> <tr> <td style="padding: 5px;">pers=third</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Person</i></td> </tr> </table>	num=sg	:	<i>Number</i>	gen=masc	:	<i>Gender</i>	pers=third	:	<i>Person</i>
phon=[man]	:	<i>Phon</i>																
cat=n	:	<i>Cat</i>																
agr	:	<table style="border-collapse: collapse; width: 100%;"> <tr> <td style="padding: 5px;">num=sg</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Number</i></td> </tr> <tr> <td style="padding: 5px;">gen=masc</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Gender</i></td> </tr> <tr> <td style="padding: 5px;">pers=third</td> <td style="padding: 5px;">:</td> <td style="padding: 5px;"><i>Person</i></td> </tr> </table>	num=sg	:	<i>Number</i>	gen=masc	:	<i>Gender</i>	pers=third	:	<i>Person</i>							
num=sg	:	<i>Number</i>																
gen=masc	:	<i>Gender</i>																
pers=third	:	<i>Person</i>																

This type identifies a unique linguistic object.  
 It is a *subtype* (or specification) of *Noun*.

## A type with several elements (underspecification of linguistic objects)

phon=[fish]	:	<i>Phon</i>									
cat=n	:	<i>Cat</i>									
agr	:	<table style="border-collapse: collapse;"> <tr> <td style="border-left: 1px solid black; border-right: 1px solid black; padding: 10px;">num</td> <td style="padding: 10px;">:</td> <td style="padding: 10px;"><i>Number</i></td> </tr> <tr> <td style="border-left: 1px solid black; border-right: 1px solid black; padding: 10px;">gen</td> <td style="padding: 10px;">:</td> <td style="padding: 10px;"><i>Gender</i></td> </tr> <tr> <td style="border-left: 1px solid black; border-right: 1px solid black; padding: 10px;">pers=third</td> <td style="padding: 10px;">:</td> <td style="padding: 10px;"><i>Person</i></td> </tr> </table>	num	:	<i>Number</i>	gen	:	<i>Gender</i>	pers=third	:	<i>Person</i>
num	:	<i>Number</i>									
gen	:	<i>Gender</i>									
pers=third	:	<i>Person</i>									

This type identifies a set of linguistic objects with phonology [fish] and category noun.

It is also a subtype (specification) of *Noun*.

## Two determiners with different degrees of specification

$$\left[ \begin{array}{l} \text{phon}=[a] : \textit{Phon} \\ \text{cat}=\text{det} : \textit{Cat} \\ \text{agr} : \left[ \begin{array}{l} \text{num}=\text{sg} : \textit{Number} \\ \text{gen} : \textit{Gender} \\ \text{pers}=\text{third} : \textit{Person} \end{array} \right] \end{array} \right]$$

$$\left[ \begin{array}{l} \text{phon}=[\text{the}] : \textit{Phon} \\ \text{cat}=\text{det} : \textit{Cat} \\ \text{agr} : \left[ \begin{array}{l} \text{num} : \textit{Number} \\ \text{gen} : \textit{Gender} \\ \text{pers}=\text{third} : \textit{Person} \end{array} \right] \end{array} \right]$$

Subtypes of *Det.*

### 3. A case where we don't seem to need unification

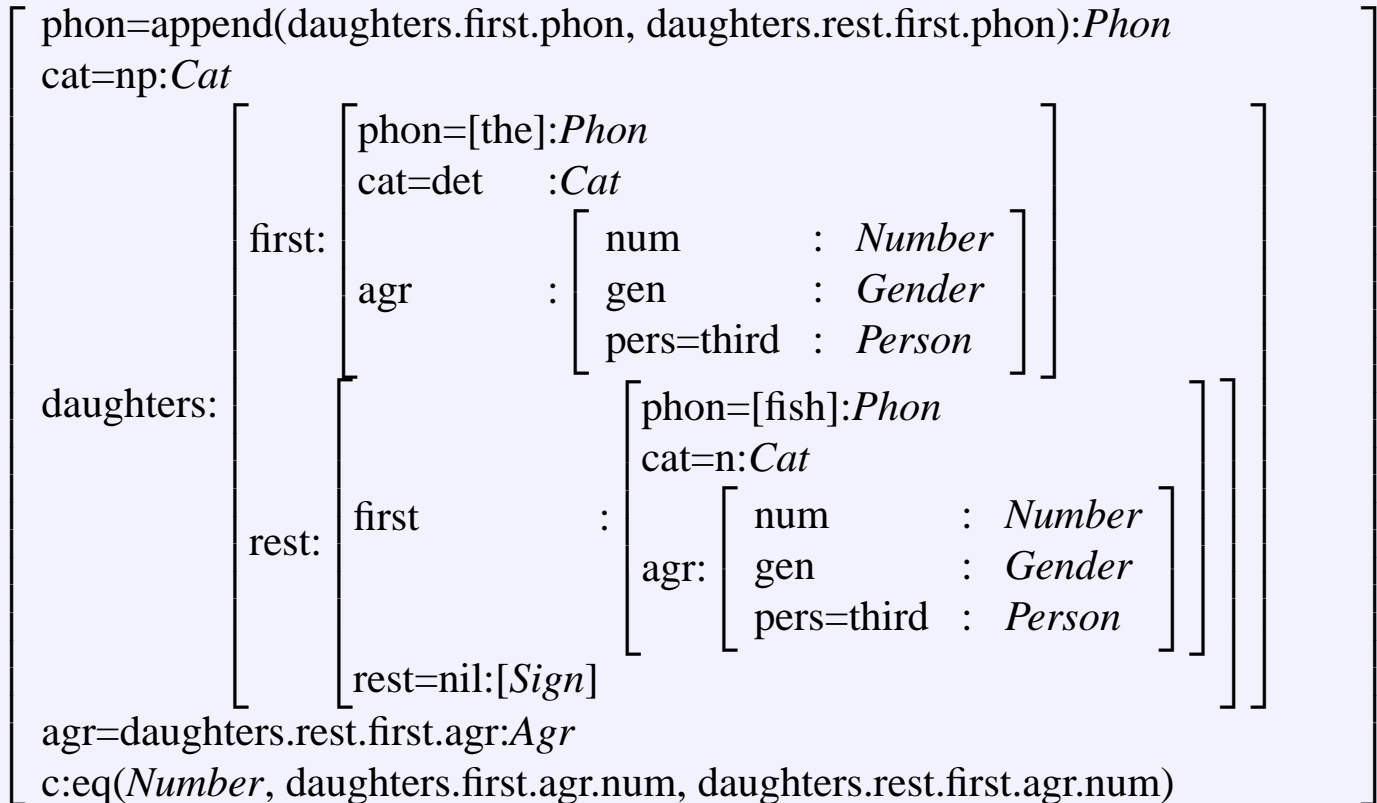
## Noun phrases (with number agreement)

$$\begin{array}{l}
 NP \equiv \\
 \left[ \begin{array}{l}
 \text{phon} = \text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}) \quad : \textit{Phon} \\
 \text{cat} = \textit{np} \quad : \textit{Cat} \\
 \text{daughters} : \left[ \begin{array}{l}
 \text{first} : \textit{Det} \\
 \text{rest} : \left[ \begin{array}{l}
 \text{first} : \textit{Noun} \\
 \text{rest} = \textit{nil} : [\textit{Sign}]
 \end{array} \right]
 \end{array} \right] \\
 \text{agr} = \text{daughters.rest.first.agr} \quad : \textit{Agr} \\
 \text{c} : \text{eq}(\textit{Number}, \text{daughters.first.agr.num}, \text{daughters.rest.first.agr.num})
 \end{array} \right.
 \end{array}$$

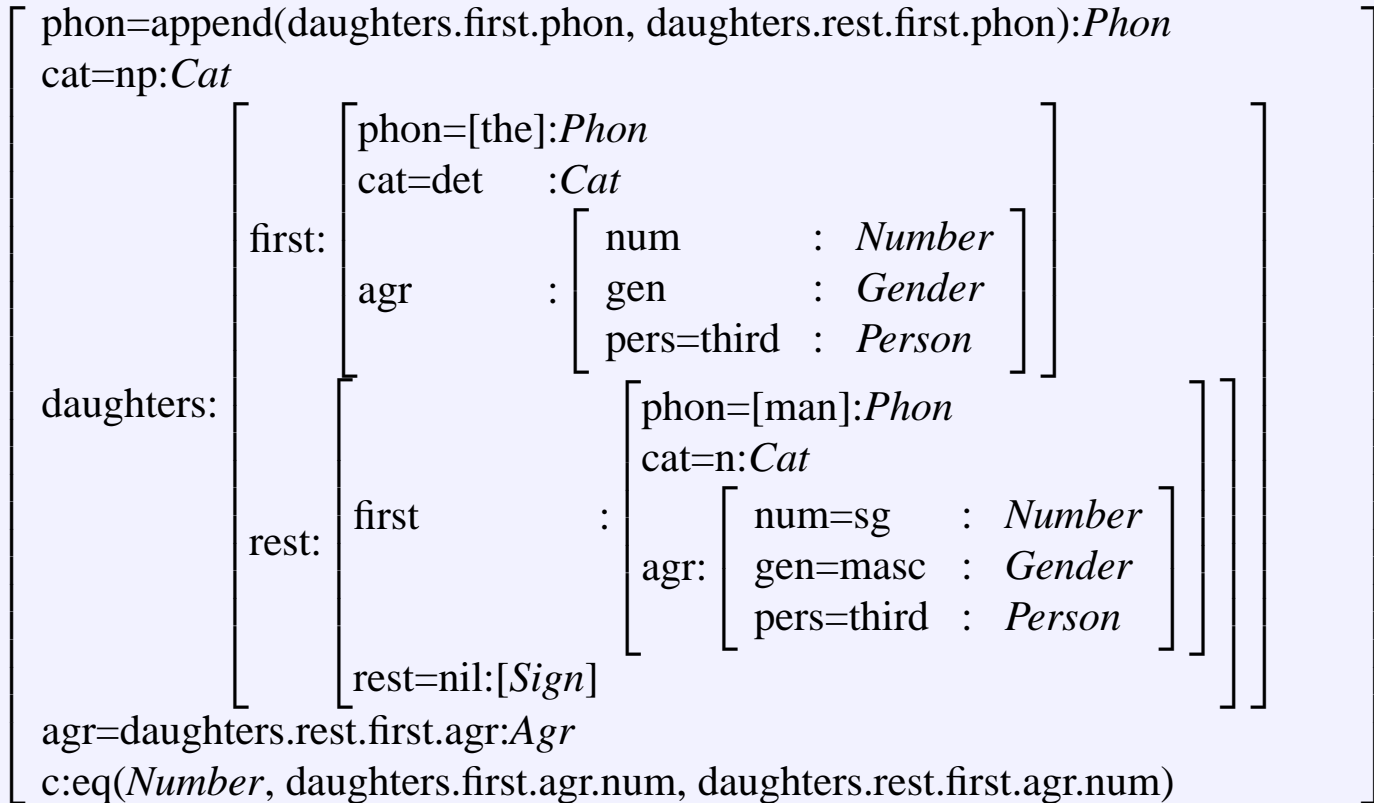
$\textit{Sign} \equiv \textit{Det} \vee \textit{Noun} \vee \textit{NP} \vee \dots$

$\text{eq}(\textit{Number}, \text{daughters.first.agr.num}, \text{daughters.rest.first.agr.num})$  is a type of proof (objects). An object of this type would be a pair of objects  $\langle a, b \rangle$  such that  $a, b : \textit{Number}$  and  $a = b$ .

# the fish

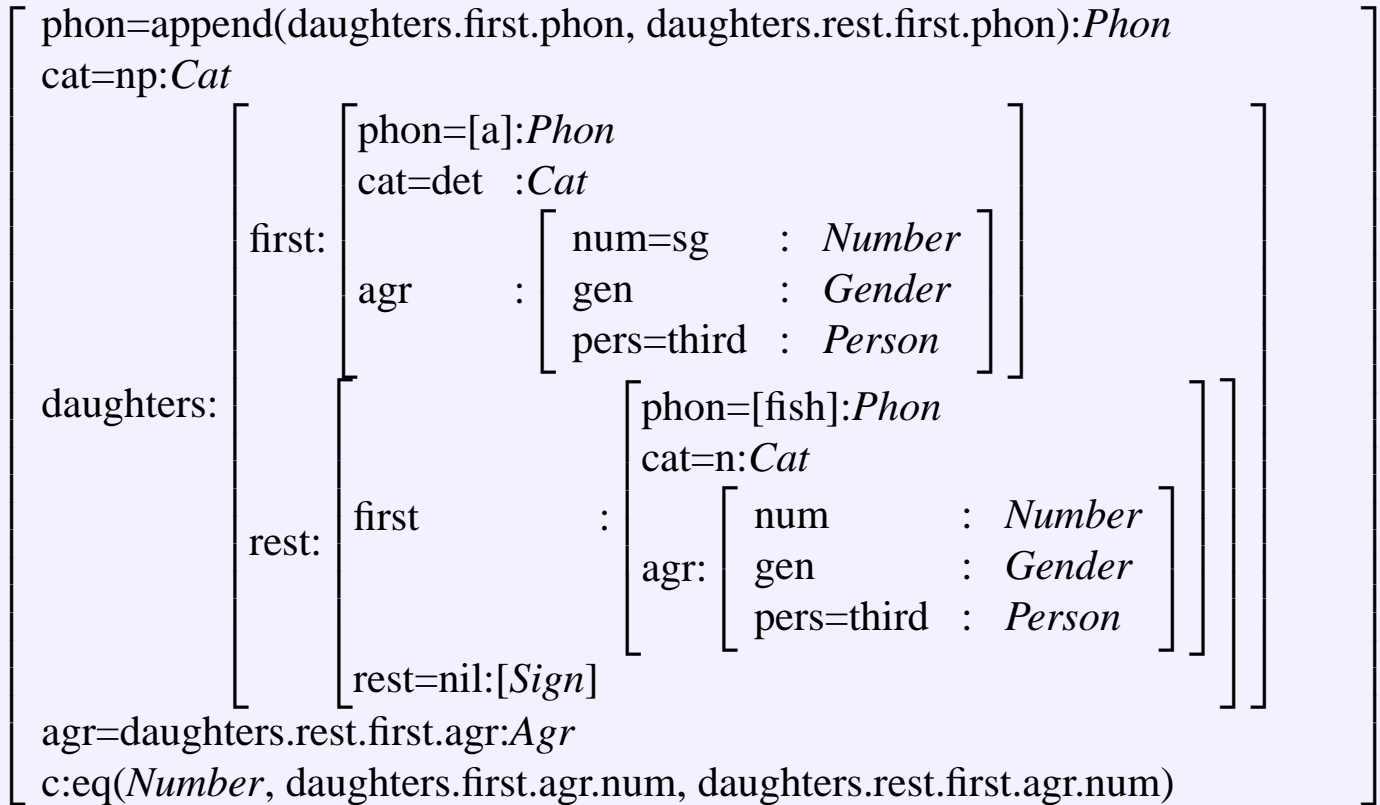


# the man



Singularity coming from *man*.

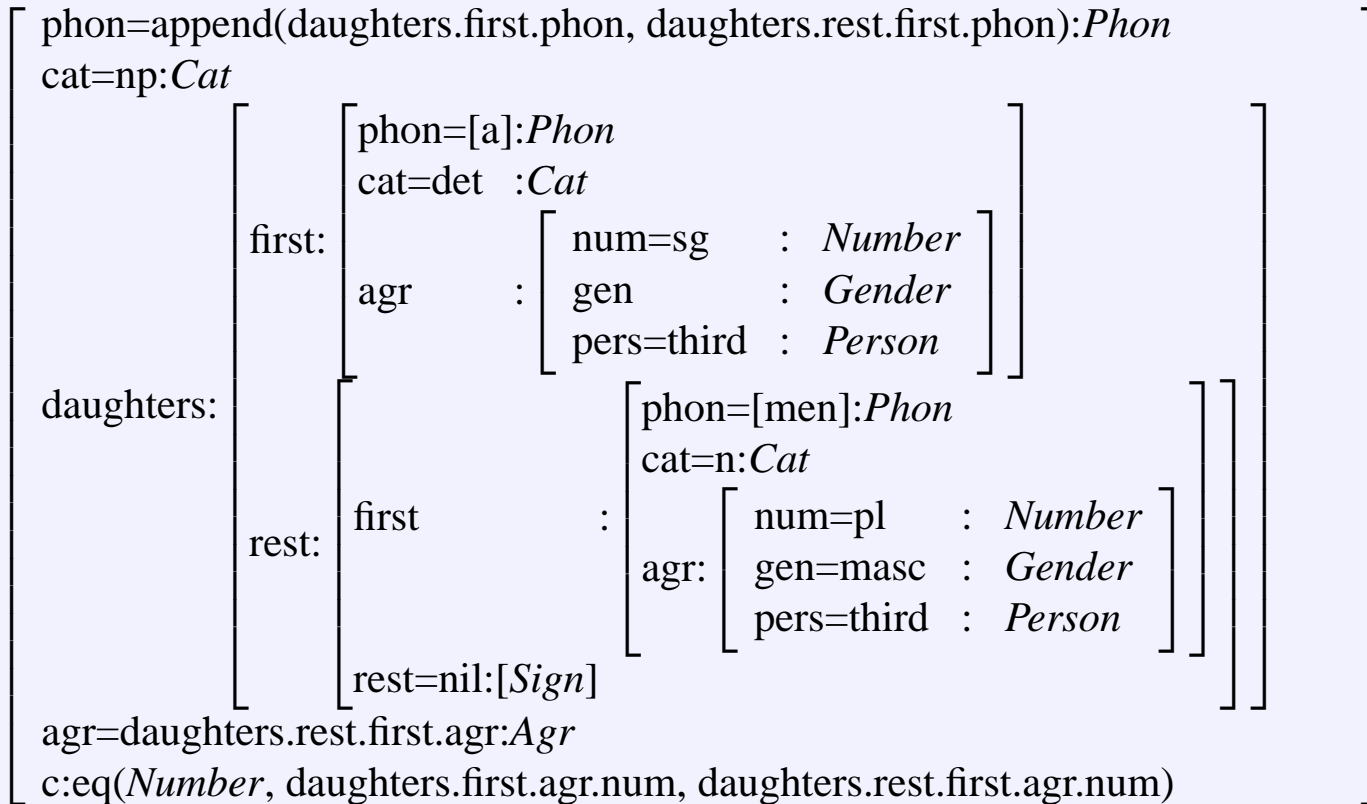
# a fish



Singularity coming from *a*.



## \*a men



An empty type – useful for robust parsing?

- types can be empty (in our case, no linguistic objects of that type)
- this does not mean that they are ill-formed types
- there are lots of distinct empty types (intensionality)
- even though a type is empty we still have a lot of information about what would have been objects of the type (contradictions are local, cf unification failure)
- e.g. if the type includes a representation of content we would still be able to interpret ungrammatical input
- potential for building a parser which assigns empty types to ungrammatical strings (in addition to types to grammatical strings)
- same intensionality that gives us an improved treatment of propositional attitudes in semantics
- therefore potentially saying something general about information and cognition, blah, blah
- cf. Frederik Fouvry's recent work on constraint relaxation with weighted feature structures (Essex PhD 2003)

## Extracting out the head feature principle

$$\begin{array}{l}
 NP \equiv \\
 \left[ \begin{array}{l}
 \text{phon} = \text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}) \\
 \text{cat} = \text{np} \\
 \text{hd-daughter} = \text{daughters.rest.first} \\
 \text{daughters}: \left[ \begin{array}{l}
 \text{first} : \textit{Det} \\
 \text{rest} : \left[ \begin{array}{l}
 \text{first} : \textit{Noun} \\
 \text{rest} = \text{nil} : [\textit{Sign}]
 \end{array} \right]
 \end{array} \right] \\
 \text{agr} \\
 \text{c} = \text{eq}(\textit{Number}, \text{daughters.first.agr.num}, \text{daughters.rest.first.agr.num})
 \end{array} \right]
 \end{array}
 \begin{array}{l}
 : \textit{Phon} \\
 : \textit{Cat} \\
 : \textit{Noun} \\
 \\
 \\
 : \textit{Agr}
 \end{array}
 \end{array}$$

^

HFP

$$\begin{array}{l}
 HFP \equiv \\
 \left[ \begin{array}{l}
 \text{hd-daughter} : \textit{Sign} \\
 \text{agr} = \text{hd-daughter.agr} : \textit{Agr}
 \end{array} \right]
 \end{array}$$

⇐ contents

## Double specification over long distances

wem Hans begegnete

(I wonder) who [dat] Hans met

$$\left[ \begin{array}{l}
 \text{phon} = \text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}) : \textit{Phon} \\
 \text{cat} = \textit{s} : \textit{Cat} \\
 \text{daughters} : \left[ \begin{array}{l}
 \text{first} : \textit{WH} \\
 \text{rest} : \left[ \begin{array}{l}
 \text{first} : \textit{S} \\
 \text{rest} = \textit{nil} : [\textit{Sign}]
 \end{array} \right]
 \end{array} \right] \\
 \text{c}_1 : \text{eq}(\textit{Cat}, \text{daughters.first.cat}, \text{daughters.rest.first.slash.cat}) \\
 \text{c}_2 : \text{eq}(\textit{Agr}, \text{daughters.first.agr}, \text{daughters.rest.first.slash.agr})
 \end{array} \right]$$

*Agr* now defined to include a case feature. *Cat* and *Agr* might be zipped together (part of local?) so that we only need one constraint.

## 4. Putting content into TTR-HPSG

# S → NP VP

$$S \equiv \left[ \begin{array}{l} \text{phon} = \text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}): \textit{Phon} \\ \text{cat} = \text{s}: \textit{Cat} \\ \text{daughters}: \left[ \begin{array}{l} \text{first}: \textit{NP} \\ \text{rest}: \left[ \begin{array}{l} \text{first}: \textit{VP} \\ \text{rest} = \text{nil}: [\textit{Sign}] \end{array} \right] \end{array} \right] \\ \text{content} = \text{daughters.first.content} @ \text{daughters.rest.first.content}: \textit{RecType} \\ \text{c}: \text{eq}(\textit{Agr}, \text{daughters.first.agr.num}, \text{daughter.rest.first.agr.num}) \end{array} \right]$$

$$\llbracket [S \text{ NP VP}] \rrbracket = \llbracket [NP] \rrbracket @ \llbracket [VP] \rrbracket$$

## VP $\rightarrow$ V NP

$$\left[ \begin{array}{l} \text{phon}=\text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}):Phon \\ \text{cat}=\text{vp}:Cat \\ \text{daughters}: \left[ \begin{array}{l} \text{first}:V \\ \text{rest}: \left[ \begin{array}{l} \text{first}:NP \\ \text{rest}=\text{nil}:[Sign] \end{array} \right] \end{array} \right] \\ \text{agr}=\text{daughters.first.agr}:Agr \\ \text{content}=\text{daughters.first.content}@\text{daughters.rest.first.content}:( [x:Ind] )RecType \end{array} \right]$$

$$\llbracket [VP \ V \ NP] \rrbracket = \llbracket V \rrbracket @ \llbracket NP \rrbracket$$

# Compositionality

*a donkey*

$$\lambda R_1:([x:Ind])\text{RecType } \lambda R_2:([x:Ind])\text{RecType} \left[ \begin{array}{l} \text{par} \quad : \quad [x : Ind] \\ \text{restr} \quad : \quad R_1 @ \text{par} \\ \text{scope} \quad : \quad R_2 @ \text{par} \end{array} \right]$$

@

$$\lambda r:([x:Ind])([c:\text{donkey}(r.x)])$$

=

$$\lambda R_2:([x:Ind])\text{RecType} \left[ \begin{array}{l} \text{par} \quad : \quad [x : Ind] \\ \text{restr} \quad : \quad [c : \text{donkey}(\text{par}.x)] \\ \text{scope} \quad : \quad R_2 @ \text{par} \end{array} \right]$$



*own a donkey*

$\lambda \mathcal{N} : (([x:Ind]) \text{RecType}) \text{RecType}$

$\lambda r_1 : [x:Ind] (\mathcal{N} @ \lambda r_2 : [x:Ind] ([c:\text{own}(r_1.x, r_2.x)]))$

@

$\lambda R_2 : ([x:Ind]) \text{RecType} \left[ \begin{array}{l} \text{par} \quad : \quad [x : Ind] \\ \text{restr} \quad : \quad [c : \text{donkey}(\text{par}.x)] \\ \text{scope} \quad : \quad R_2 @ \text{par} \end{array} \right]$

=

$\lambda r_1 : [x:Ind] \left( \left[ \begin{array}{l} \text{par} \quad : \quad [x : Ind] \\ \text{restr} \quad : \quad [c : \text{donkey}(\text{par}.x)] \\ \text{scope} \quad : \quad [c : \text{own}(r_1.x, \text{par}.x)] \end{array} \right] \right)$

## 5. A case where we seem to need not to have unification

# Information structure

Vallduví, Engdahl and Vallduví

- [The president hates the Delft china set <sub>F</sub>]
- The president [hates the Delft china set <sub>F</sub>]
- The president hates [the Delft china set <sub>F</sub>]
- The president [hates <sub>F</sub>] the Delft china set

## The problem

In a classical HPSG approach using unification the content of the verb is the same as the content of the sentence

# The syntactic strategy

Engdahl and Vallduví

Mark focus on syntactic constituents

Still leaves us with no representation of focus information

## The MRS strategy

Wilcock

every dog chased [some cat <sub>F</sub>]

1:every(x,3,5), 3:dog(x), 7:cat(y), 5:some(y,7,4), 4:chase(e,x,y)

TOP-HANDLE:1, LINK:1, FOCUS:5

You seem to get different focus information depending on how you plug things together? Notice that scope and information structure are independent.

## The separate semantic formalism strategy

Ericsson

From the unification formalism point of view semantic representations (e.g. using the  $\lambda$ -calculus) are atoms

This works fine because no unification is going on in the semantic representation.

But it misses the integration of syntax and semantics that HPSG wants to have.

## Towards a TTR-HPSG analysis

We show some expressive possibilities, though not yet an general analysis.



# VP $\rightarrow$ V [F NP]

```

[ phon=append(daughters.first.phon, focus(daughters.rest.first.phon)):Phon
  cat=vp:Cat
  daughters: [ first:VSign
               rest: [ first:NPSign
                       rest=nil:[Sign] ] ]
  content=daughters.first.content@daughters.rest.first.content:RecType
  info-struct: [ focus=daughters.first.rest.first.content:Quant
                 ground=daughters.first.content:(Quant)Property ] ]
  
```

```

[phon=append(daughters.first.phon, daughters.rest.first.phon):Phon
cat=s:Cat
daughters: [first:NPSign
            rest: [first:VPSign
                  rest=nil:[Sign]]]
content=daughters.first.content@daughters.rest.first.content:RecType
info-struct: [focus=daughters.first.rest.first.focus:Quant
              ground= $\lambda q:Quant$ 
                  (daughters.first.content
                   @(daughters.rest.first.ground@q)):RecType]

```

## 6. Conclusions

Type theory with records provides us with a powerful formalism for

- feature-based analyses
- “unification” phenomena
- semantics
- “functional” phenomena

and, importantly, an *integrated* approach.

## A. Abstract and concrete syntax

## Names for abstract and concrete syntax

abstract syntax	concrete syntax	Ranta GF, compiler technology
tectogrammar	phenogrammar	Curry, Dowty
categorial grammar derivation	strings at nodes in derivation	Montague's CG based grammar

## Properties of systems which distinguish abstract and concrete syntax

- Concrete syntax is a *compositional* projection from abstract syntax
- Abstract syntax can be a semantic representation, but may be a more abstract syntactic representation (cf Montague Grammar)
- A grammar can be regarded as a function from expressions of abstract syntax to sets of expressions of concrete syntax (Ranta)
- There may be several concrete syntaxes corresponding to a single abstract syntax (Ranta)
- Grammars (functions) can be composed (Ranta)

## Related compositional notions

Abstract	Concrete
CF-backbone	HPSG feature structure
C-structure	F-structure
C-structure	Phonology projection

## An intuitively related non-compositional notion

Abstract	Concrete
D-structure	S-structure
D-structure	LF
D-structure	P-structure



# Extracting the categorial abstract syntax from HPSG signs

Use the `cat` and `daughters` attributes.

$$\llbracket [S \text{ NP VP}] \rrbracket = \left[ \begin{array}{l} \text{phon} = \text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}): \textit{Phon} \\ \text{cat} = \textit{s:Cat} \\ \text{daughters:} \left[ \begin{array}{l} \text{first} = \llbracket \text{NP} \rrbracket : \textit{NP\textit{Sign}} \\ \text{rest:} \left[ \begin{array}{l} \text{first} = \llbracket \text{VP} \rrbracket : \textit{VP\textit{Sign}} \\ \text{rest} = \textit{nil:Sign} \end{array} \right] \end{array} \right] \\ \text{content} = \text{daughters.first.content} @ \text{daughters.rest.first.content}: \textit{RecType} \end{array} \right]$$

⇐ contents

## Using a more GF-like notation

```

predS np vp =
[
  phon=append(daughters.first.phon, daughters.rest.first.phon):Phon
  cat=s:Cat
  daughters: [
    first=np:NPSign
    rest: [
      first=vp:VPSign
      rest=nil:[Sign]
    ]
  ]
  content=daughters.first.content@daughters.rest.first.content:RecType
]

```

# Extracting function argument abstract syntax from HPSG signs

Use the content attributes.

apply np vp =

$$\left[ \begin{array}{l} \text{phon}=\text{append}(\text{daughters.first.phon}, \text{daughters.rest.first.phon}):Phon \\ \text{cat}=\text{s}:Cat \\ \text{daughters}: \left[ \begin{array}{l} \text{first}=\text{np}:NP\text{Sign} \\ \text{rest}: \left[ \begin{array}{l} \text{first}=\text{vp}:VP\text{Sign} \\ \text{rest}=\text{nil}:[Sign] \end{array} \right] \end{array} \right] \\ \text{content}=\text{daughters.first.content}@ \text{daughters.rest.first.content}:RecType \end{array} \right]$$

⇐ contents

## Projecting the content and phonology

`(apply np vp).content = np.content@vp.content`

`(apply np vp).phon = append(np.phon, vp.phon)`

## HPSG and LFG

- HPSG provides data structures in which abstract syntax is hidden (but can be extracted)
- LFG uses C-structure as abstract syntax and projects to different representations (which could be zipped together in a single data structure)
- A true grammatical framework will allow us to convert grammars between these and other possibilities (e.g. alternative abstract syntaxes)
- It will also allow high level reasoning about grammars (alternative concrete syntaxes, grammar composition, grammar extraction)